

HEURISTIC BASED TASK SCHEDULING IN MULTIPROCESSOR SYSTEMS WITH GENETIC ALGORITHM BY CHOOSING THE ELIGIBLE PROCESSOR

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ABSTRACT

In multiprocessor systems, one of the main factors of systems' performance is task scheduling. The well the task be distributed among the processors the well be the performance. Again finding the optimal solution of scheduling the tasks into the processors is NP-complete, that is, it will take a lot of time to find the optimal solution. Many evolutionary algorithms (e.g. Genetic Algorithm, Simulated annealing) are used to reach the near optimal solution in linear time. In this paper we propose a heuristic for genetic algorithm based task scheduling in multiprocessor systems by choosing the eligible processor on educated guess. From comparison it is found that this new heuristic based GA takes less computation time to reach the suboptimal solution.

KEYWORDS

Multiprocessor, task scheduling, heuristic, genetic algorithm

1. INTRODUCTION

In multiprocessor based system the processing capability of processors may vary. The parallel tasks must be allocated into the processors such that the total completion time must be as less as possible. The optimal usage of processors is also expected. Again optimal task scheduling in multiprocessor systems is NP-complete [8], that is, finding optimal scheduling of tasks for multiprocessors is time consuming. These define the problem of task scheduling on multiprocessor systems to allocate a set of tasks to processors such that the optimal usage of processors and accepted computational time for scheduling algorithm are obtained. Genetic Algorithm, an evolutionary algorithm is used to find a suboptimal solution of the problem in considerable computation time. To reach the solution faster many heuristic based approach are used. By heuristic based approach the initial population is much closer to optimal solution. This results much less computation time in GA. In this paper, we propose a modification of heuristic approach of genetic algorithm method based on bottom-level by choosing the eligible processor for assigning the tasks which eventually decreases the computational time for finding the suboptimal schedule.

This paper is based on deterministic model, that is task dependencies and their execution time are known. The communication costs among the tasks are negligible and the numbers of

processors are also fixed. The dependencies along with execution time of the tasks are represented by a Directed Acyclic Graph (DAG).

The remainder of this paper is organized as follows: brief summary of the related works, explanation of DAG in section 3, Heuristic based Genetic Algorithm explanation in section 4, the proposed improvement in section 5, the result of experimental studies are presented on section 6 and a conclusion and future work in section 7.

2. RELATED WORK

Since the beginning of the research on this field many approaches have been developed to solve the task scheduling on multiprocessor system. Some are heuristic based approach [9-11]; some rely on evolutionary approaches [12-14] and some follows the hybrid methods [15-18].

There are many heuristic based methods for solving multiprocessor task scheduling approach. The best heuristic approaches are based on task list technique [19]. In this technique, a list of descending priority ordered tasks is made. A task is selected from the head of the list and assigned to the processor. This method is also classified as static and dynamic. In static scheduling algorithms the list is not updated with new ordering at run time while the dynamic approaches do. Scheduling algorithms using t-level (top level) and b-level (bottom level) attributes for assigning priority to the processors have been proposed. There is another frequently used parameter ALAP (As Late As Possible) start time [20], which defines the longest possible execution time that a task can be postponed. List scheduling ISH (Insertion Scheduling) followed by DSH (Duplication Scheduling) that is a task duplication method has been also proposed [21]. There are several other Heuristic methods (Level-based Heuristics) [22] such as HLFET ((Highest Level First with Estimated Times), HLFNET (Highest Levels First with No Estimated Times), Random (the assigned tasks priority are random), SCFET (Smallest Co-levels First with Estimated Times) and SCFNET (Smallest Co-levels First with No Estimated Times), CP/MISF (critical path/most immediate successors first) [10], HNF (Heavy Node First) and WL(Weighted Length) [19]. All of these attributes act based on level concept in the DAG and without consideration of communication cost. Moreover, DF/IHS [10], EZ (Edge-zeroing) algorithm [23], LC (Linear Clustering) algorithm [24], DSC (Dominant Sequence Clustering) algorithm [25], MD (Mobility Directed) [27], DCP (Dynamic Critical Path) [11], ETF (Earliest Task First) [9] and greedy heuristics [26] are other heuristic methods.

These heuristic based methods are not considered as intense as before as they do not have good result in all cases. Therefore research on combinatorial optimization algorithms such as GA [2, 3,12, 13, 28], meta-heuristics and hybrid methods [5, 18] are going on.

3. THE DAG MODEL

A Directed Acyclic Graph (DAG) for tasks is the graph that represents the precedence constraints among the tasks along with their execution time. The DAG can be represented by a set $G = \{V,E\}$ where V is the set of the task and E is the set of relations between the tasks. When the DAG is represented in graph, V represents the nodes and E represents the edges among the nodes. The computation cost of a task is represented the by the weight of the node and is denoted by $W(T_i)$ where T_i is the i th task. Figure 1 represents a DAG of 11 tasks along with their precedence constraints. Each edge in the DAG represents the relationship between the tasks.[23]

If there is an edge E_{ij} from task T_i to task T_j then task T_i precedes the task T_j that is T_i is the predecessor of T_j and T_j is successor of T_i . It is also represented by the $T_i \Rightarrow T_j$. The height of a task height (T_i) can be represented by

$$height(T_i) = \begin{cases} 0 & \text{if } PRED(T_i) = \emptyset \\ 1 + \max height(T_j) | T_j \in PRED(T_i) & \text{otherwise (1)} \end{cases}$$

where $PRED(T_i)$ is the set of predecessors of the task T_i .

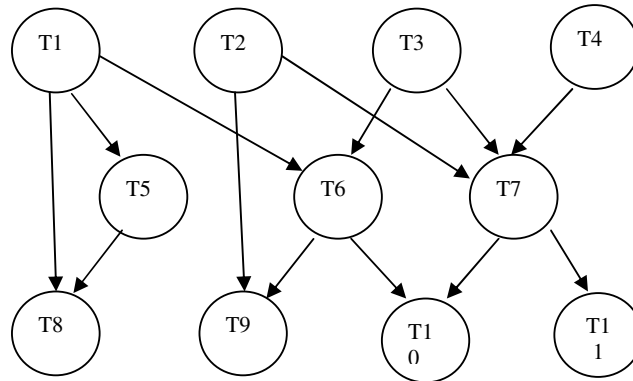


Figure 1. An example of a DAG

If a task T_i is assigned to a processor then $ST(T_i)$ and $FT(T_i)$ denote the start time and finishing time respectively. When all the tasks are scheduled, $\text{MAX}\{FT(T_i)\}$ denotes the schedule length across all processors. The goal of scheduling is to minimize the $\text{MAX}\{FT(T_i)\}$.

In this paper we consider a set of processors P where all the processors are homogeneous i.e. all the processor will have same execution time to run a task individually. The number of processors is bounded and all the processors are connected with negligible communication cost.

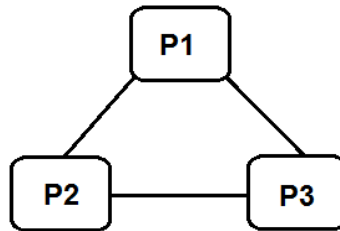


Figure 2 represents three fully connected homogeneous parallel systems.

Figure 2. Three fully connected machines

4. GENETIC ALGORITHM

A genetic algorithm is an evolutionary algorithm which generates near optimal solution of a problem by a guided random search method where elements (called individuals) in a given set of solutions (called population) are randomly combined and modified until some termination condition is achieved. The population evolves iteratively in order to improve a given cost function or fitness function of its individual [4]. In our case, the individuals are all the task-

processor pairs and combinations of multiple individuals form the term population. The fitness of a population is $\text{MAX} \{FT (P_i)\}$ for $j = 1, 2, 3, \dots, P_n$, where P_n is the number of processors and FT is the finishing time of the final task in processor P_i . The objective of GA is to find the minimum $\text{MAX} \{FT (P_i)\}$.

The initial step for a genetic algorithm is to get a set of initial population. Typically these populations are generated randomly. Then these populations are reviewed based on fitness function. Good individuals are replicated while bad individuals are removed. After selecting a pair of parent solution a crossover operation is performed to produce child solutions which preserve the characteristic of their parents. The mutation operation is randomly performed so that the random search algorithm does not stick in local minima. The mutation operation injects new characteristics in population to explore the uncovered area of the random search.

5. SCHEDULING ALGORITHMS

5.1. Heuristic Based Scheduling Algorithms

For the bounded number of processors (BNP) there are several scheduling algorithms i.e. HLFET algorithm, ISH algorithm, MCP algorithm, ETF algorithm, DLS algorithm. From them HLFET [22] is the simplest algorithms based on list scheduling technique. The basic idea of list scheduling is to make a scheduling list (a sequence of nodes for scheduling) by assigning them some priorities, and then repeatedly execute the following two steps until all the nodes in the graph are scheduled:

- (1) Remove the first node from the scheduling list;
- (2) Allocate the node to a processor which allows the earliest start-time.

There are various ways to determine the priorities of nodes. HLFET algorithm uses b-level based priority scheduling. The HLFET algorithm can be described as below

- (1) Calculate the static b-level (i.e., sl or static level) of each node.
- (2) Make a ready list in a descending order of static b-level. Initially, the ready list contains only the entry nodes. Ties are broken randomly.

Repeat

- (3) Schedule the first node in the ready list to a processor that allows the earliest execution, using the non-insertion approach.
- (4) Update the ready list by inserting the nodes that are now ready.

Until all nodes are scheduled.

5.2. Heuristic Based Genetic Algorithm

Performance of GA greatly depends on initial population as the more fit the initial population the faster it converges towards suboptimal solution.

Each individual of the initial population is generated through a minimum execution time or min-min heuristic along with b-level or t-level precedence resolution to avoid the problem of same execution time or completion time and same precedence. The problem of same execution time/completion time and precedence can occur in the homogeneous parallel system as all the processors take same execution time to execute one task.

The task to be scheduled for each iteration using b-level precedence resolution is determined by the following rules:[1,6]

- (1) Sort the tasks according to their heights in ascending order.
- (2) Sort the tasks with the same height according to their bottom-level in descending order.
- (3) Repeat step 4 and step 5 until finish of all the tasks.
- (4) Generate a permutation of processors.
- (5) Assign tasks to processors in order.
- (6) The above steps are repeated for the number of population size.

The length of all individuals in an initial population is equal to the number of tasks in the DAG. Following Table 1 represents the execution time and the priority of tasks' execution based on their bottom-level of the DAG presented in Figure 1.

Task number	Execution time	Height	Bottom-level	Order of execution according to bottom-level
1	50	0	72	1
2	1	0	41	4
3	10	0	50	3
4	20	0	60	2
5	20	1	21	7
6	2	1	22	6
7	20	1	40	5
8	1	2	1	11
9	20	2	20	9
10	19	2	19	10
11	20	2	20	8

Table 1: Execution time and priority of execution based on bottom-level

Figure 3 shows schedule for three processors based on the order of execution according to the bottom-level. From the figure we see that the total finish time based on the priority of the tasks' bottom-level is 92.

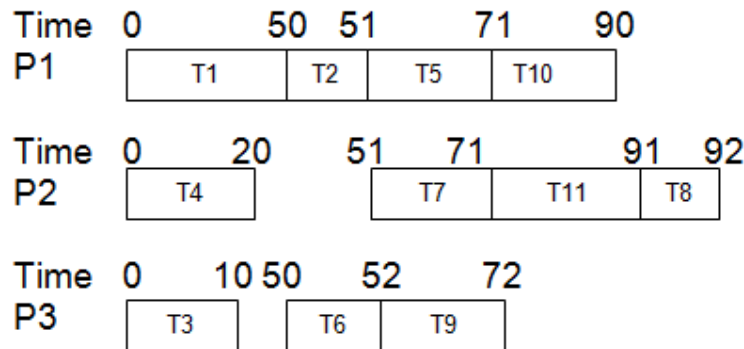


Figure 3: Schedule for three processors based on the Order of execution according to the bottom-level

5.3. Variation of HLFET heuristic over GA

In step 4 of Heuristic based GA using b-level precedence resolution a permutation of processors is performed. Instead of choosing processor by performing permutation of processors we propose the following algorithm.

Let the task mapping on the processors be represented by $A\{T, P\}$ where $T = \{T_1, T_2, \dots, T_n\}$, a set of n tasks and $P = \{P_1, P_2, \dots, P_m\}$, a set of m processors. $A\{T, P\}$ represents that task $T_i \in T$ is mapped on processor $P_k \in P$. A set of unassigned tasks is represented as $U = \{T_1, T_2, \dots, T_n\}$ where the tasks are sorted in descending order by b-level precedence. Tasks weight T_w is the execution time of task T . Processor's weight P_{kw} is the total execution time and waiting time of the processor so far.

$$P_{k_w} = \sum_{i=0}^n T_{i_w} + W_k$$

where T_i is assigned task on processor P_k and T_{i_w} is the execution time of task T_i and W_k is the total waiting time of the processor P_k . $P_k T_{i_w}$ is the processors weight when the task T_i is assigned on processor P_k . Procedure $LastTask(P_k)$ returns the last assigned task on processor P_k .

The pseudo code of our algorithm is:

```

Begin:
  While( $U \neq \phi$ )
    Begin:
      Select first task  $T_j$  from unassigned task set  $U$  and remove the task from  $U$ .
      Find  $P_k$  from  $\max(P_k T_{i_w}, T_i \in PRED(T_j))$ .
      If( $P_k \neq \phi$  &&  $T_i = LastTask(P_k)$ )
        Begin:
          Assign  $A\{T_j, P_k\}$ 
           $P_{k_w} += T_{j_w}$ 
        End
      Else If( $P_k \neq \phi$ )
        Begin:
          Find  $P_k = \min(P_{k_w})$ 
          Assign  $A\{T_j, P_k\}$ 
           $P_{k_w} = \max(P_k T_{i_w}, T_i \in PRED(T_j)) + T_{j_w}$ 
        End
      Else
        Begin:
          Find  $P_k = \min(P_{k_w})$ 
          Assign  $A\{T_j, P_k\}$ 
           $P_{k_w} += T_{j_w}$ 
        End
      End If
    End of While
  End

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Table 2: Variation of HLFET heuristic over GA

By applying this algorithm on the DAG of figure 3, we get the schedule for three processors system of figure 4.

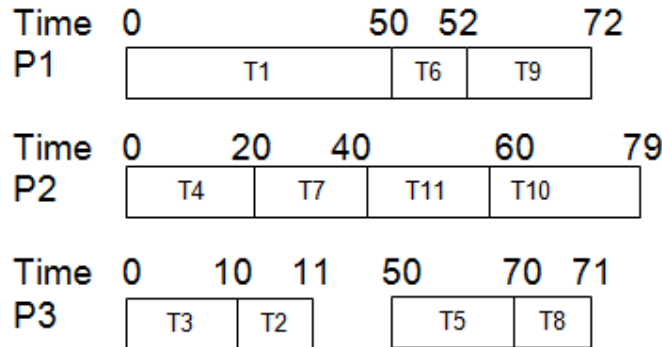


Figure 4: Schedule for three processors based on our proposed variation of HLFET heuristic over GA

From the figure we see that the total finish time of our algorithm is 79 where as the total finish time of round robin processor selection is 92.

6. EXPERIMENT RESULTS

Standard Task Graph (STG) has been used to benchmark the evaluation of multiprocessor scheduling algorithms [7]. All the task graphs in STG had been generated randomly without communication cost. As it is hard to evaluate on all the task graphs, we have chosen several task graphs randomly from the standard STG and performed the evaluation of the Elitism algorithm and our proposed scheduling algorithm. The tests have been run on 3, 4 and 16 processors with 50 and 100 tasks. For each set of test the test has been run for 10 times and from them average makespan - best makespan, average and best number of evaluations (to reach the termination condition) have been found.

Table 3 shows the simulation result of Elitism algorithm and our proposed algorithm in terms of average makespan, Best finish time, average and best number of evaluations to reach the termination condition.

Problem	Elitism algorithm				Proposed Algorithm			
	Avg costs:	Best costs:	Avg evaluation	Best evaluation	Avg costs:	Best costs:	Avg evaluation	Best evaluation
50tasks\Rand0100\4processors	0.00084864	0.0008336	182.8	150	0.00085296	0.0008256	169.3	135
50tasks\Rand0100\16processors	0.00154016	0.0013208	58.7	1	0.00149448	0.0011544	78.1	1
50tasks\Rand0069\4processors	0.00212304	0.0020992	146.5	96	0.0019688	0.0019688	1	1
50tasks\Rand0069\16processors	0.00226752	0.0022632	20.5	1	0.0019688	0.0019688	1	1
50tasks\Rand0019\4processors	0.00071768	0.000704	158.9	131	0.00066264	0.000656	87.9	12
50tasks\Rand0019\16processors	0.00088928	0.00078	190.3	165	0.00079704	0.0007704	158.3	81
50tasks\Rand0016\4processors	0.00069064	0.0006736	151.9	108	0.00064984	0.0006392	58.4	1
50tasks\Rand0016\16processors	0.00082856	0.0007992	173.8	128	0.00076744	0.000704	180.2	135
50tasks\Rand0002\4processors	0.00050512	0.0004744	169	106	0.00049024	0.0004712	162.3	119
50tasks\Rand0002\3processors	0.0000696	0.0000685	62.7	10	0.0000694	0.000069	46.7	5

50tasks\Rand0002 \16processors	0.0008328	0.0006888	190	171	0.00060336	0.000548	192.4	176
100tasks\Rand010 0\4processors	0.00121952	0.0011928	179.9	140	0.00122984	0.0011384	171.1	136
100tasks\Rand010 0\16processors	0.002604	0.002604	1	1	0.002604	0.002604	1	1
100tasks\Rand006 9\4processors	0.001816	0.0017928	147.1	68	0.00179592	0.001776	143.2	53
100tasks\Rand006 9\16processors	0.00355936	0.0033864	181.3	151	0.0031724	0.0029736	182.2	155
100tasks\Rand001 9\4processors	0.00205848	0.002024	147	64	0.00203688	0.002016	134	75
100tasks\Rand001 9\16processors	0.00429416	0.0039576	154.1	1	0.00386936	0.0034488	170.4	91
100tasks\Rand001 6\4processors	0.00190968	0.0018432	174.9	101	0.00185688	0.001784	159.4	122
100tasks\Rand001 6\16processors	0.00410888	0.0038808	164.8	1	0.00357968	0.0032568	171.6	97
100tasks\Rand000 2\4processors	0.00169768	0.0015488	171.9	139	0.00167768	0.0015448	169.3	140
100tasks\Rand000 2\16processors	0.002696	0.002696	1	1	0.002696	0.002696	1	1

Table 3: comparison between Elitism algorithm and our proposed algorithm

Figure 5 represents the comparison of average number of evaluations to reach the termination condition with Elitism Algorithm and our proposed algorithm. From the figure it is clear that our proposed algorithm takes much less time to reach the termination condition than Elitism Algorithm except the cases where there are 16 processors. While evaluating the figure 6 which represents the comparison of average make span we find that our algorithm performs much better than Elitism algorithm even for the 16 processors.

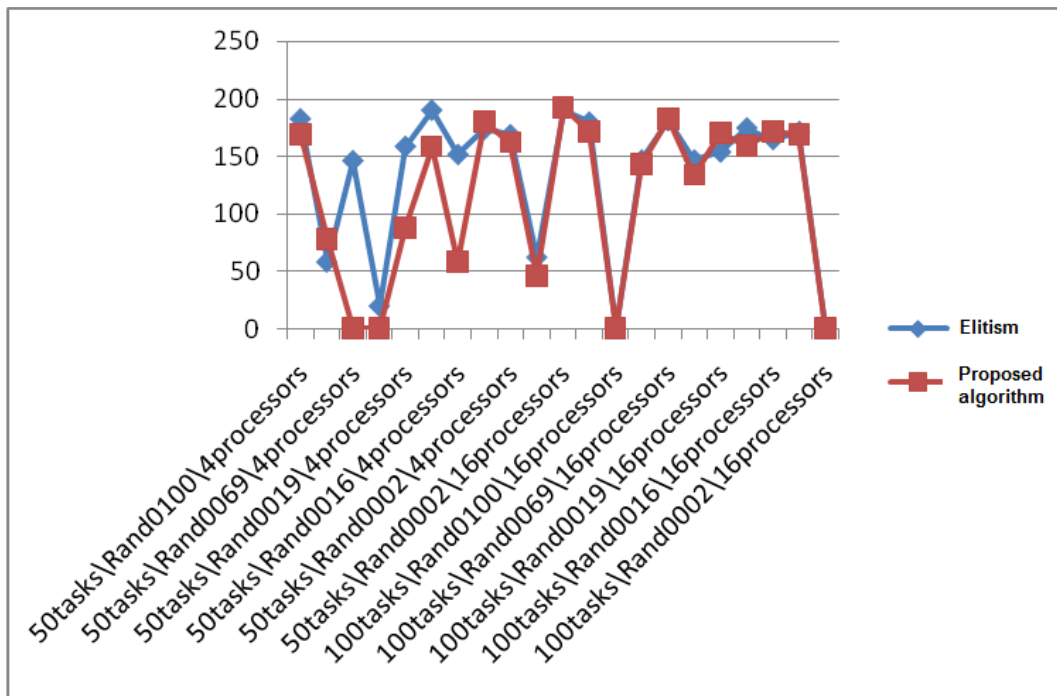


Figure 5: comparison of average number of evaluations to reach the termination condition with Elitism Algorithm and our proposed algorithm

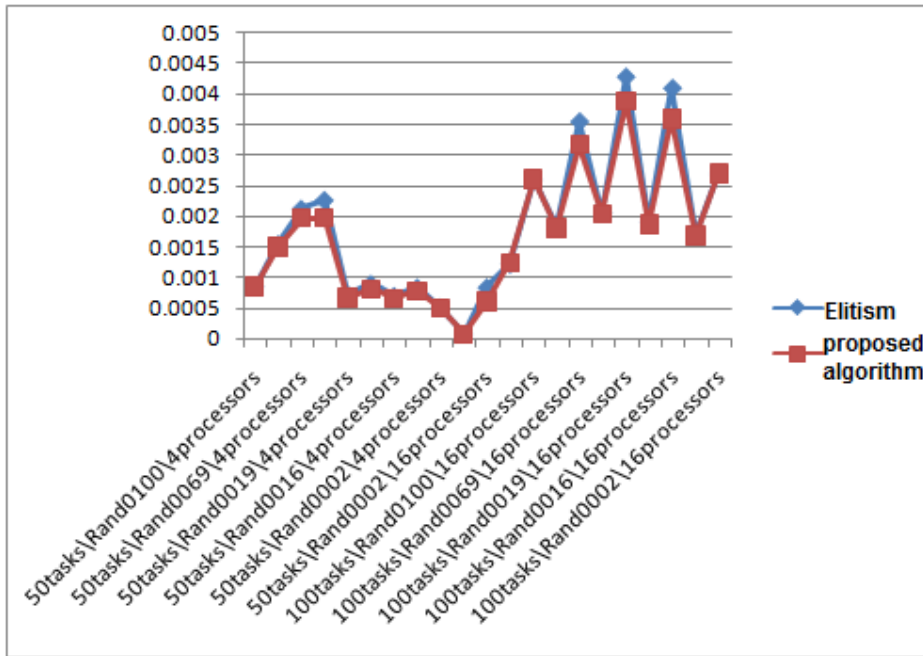


Figure 6: comparison of average makespan with Elitism Algorithm and our proposed algorithm

Figure 7 shows the steps of evaluation of a test case. In the test case there are 100 task with 16 processors. The test data is based on Rand069. The evaluation steps are represented along x-axis and the maximum makespan or cost is represented on y-axis. From the plot it is visible that the evaluation path of our algorithm is beneath of the Elitism algorithm. This means that our algorithm schedules the tasks with lower makespan quickly than the Elitism Algorithm.

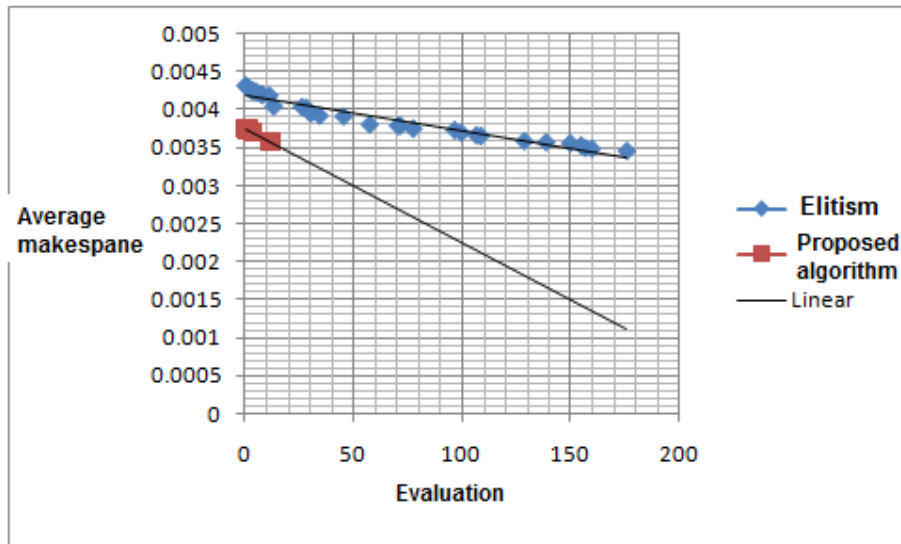


Figure 7: comparison of average makespan - evaluation with Elitism Algorithm and our proposed algorithm for 100 task with 16 processors (Rand069)

7. CONCLUSIONS

In this paper we have discussed the HLFET algorithm and Elitism Stepping method for task scheduling in multiprocessor systems with genetic algorithm. We also proposed a variation of HLFET algorithm and simulated the algorithm as well as Elitism stepping method with Standard Task Graph (STG). We have compared the simulation result and found that our proposed algorithm has better average makespan in smaller number of evaluations than Elitism Stepping method.

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